

# FIXED-PRIORITY SCHEDULING OF DUAL-CRITICALITY SYSTEMS

Sanjoy Baruah and Alan Burns

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1. **CONTEXT** AND BACKGROUND
2. FP SCHEDULING AND MIXED CRITICALITIES: A **SURVEY**
3. **PERIOD TRANSFORMATION** (PT) AND MIXED CRITICALITIES

# CONTEXT AND BACKGROUND

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## Requirements in real-time systems design

**Correctness** - by using **rigorous design methodologies**

Constrain the system designer to only choose **safe** designs

“Freedom from choice” - A. Sangiovanni-Vincentelli

**Efficiency** - by the application of **scheduling theory**

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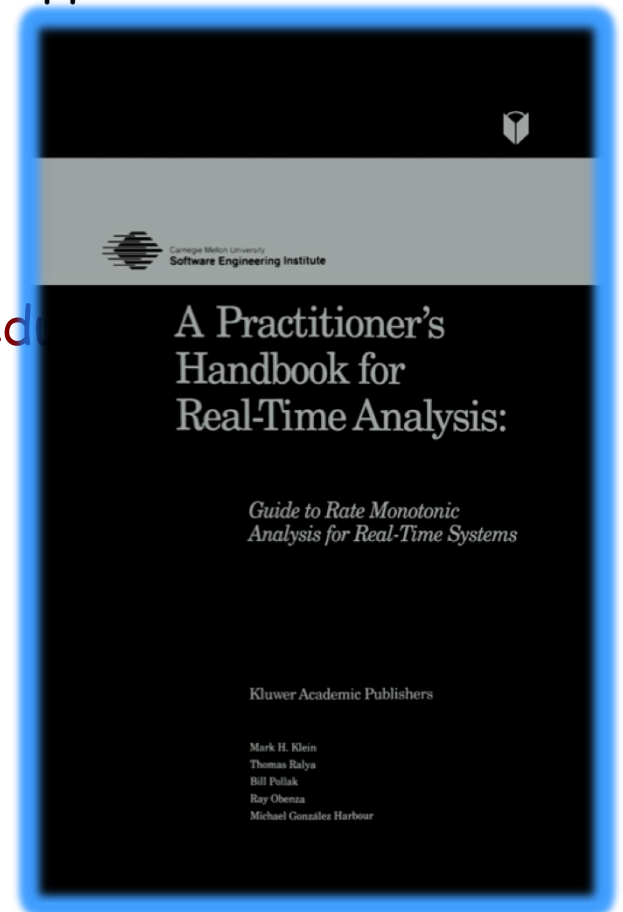
## Requirements in real-time systems design

**Correctness** - by using **rigorous design methodologies**

Fixed Priority (FP) scheduling: the **RMA** approach

- Formalized industrial practice
- Based on strong mathematical foundations
- Good out-reach (and public relations! 😊)

**Efficiency** - by the application of **sched**



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## Requirements in real-time systems design

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Fixed Priority (FP) scheduling: the **RMA** approach

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**Efficiency** - by the application of **scheduling theory**

Increase efficiency:

“Better” **priority assignment**: **Rate/ Deadline Monotonic**

Choosing the **periods** appropriately: **harmonic** task systems

# A complicating factor

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RT systems becoming more **complex** ➡ implemented upon **COTS** platforms

COTS platforms are **nondeterministic** at run-time

Example:  **$x := a + b$**  on the Motorola PowerPC 755

- Best case: 3 cycles

- Worst case: 321 cycles

In the 1990's: the Motorola 64K had best case = worst case = 20 cycles

Resource over-provisioning to deal with unpredictability

Resource over-provisioning causes **run-time inefficiencies**

# Mixed-criticality systems

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Steve Vestal. *Preemptive scheduling of multi-criticality systems with varying degrees of execution time assurance*. [RTSS 2007](#)

Functionalities of **different criticalities** upon a shared platform

- use **resource over-provisioning** to validate **highly critical** functionalities at **high levels of assurance**

- **design-time** resource reclamation to validate **less critical** functionalities at **lower levels of assurance**

**A NEW scheduling problem (finally! 😊)**



# **FIXED-PRIORITY SCHEDULING OF MIXED-CRITICALITY SYSTEMS: A REVIEW**

# Priority assignment schemes

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Collection of mixed-criticality **sporadic tasks** executing on a **single preemptive processor**

Each task  $\tau_i$ :

- **HI** or **LO** criticality
- Relative deadline  $D_i$
- Minimum inter-arrival separation (“period”)  $T_i$
- Two WCETs:  $C_i(\text{LO})$  and  $C_i(\text{HI})$ , with  $C_i(\text{LO}) \leq C_i(\text{HI})$

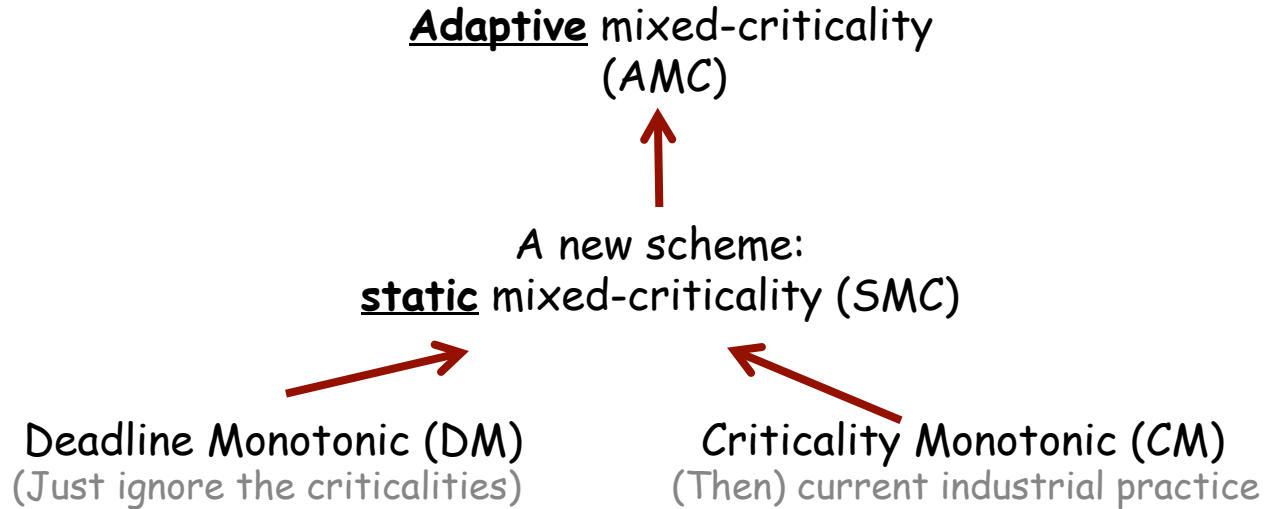
**Determine** a single scheduling strategy

- If each job of each  $\tau_i$  completes in  $\leq C_i(\text{LO})$  units of execution, then all jobs of all tasks complete by their deadlines
- If each job of each  $\tau_i$  completes in  $\leq C_i(\text{HI})$  units of execution, then **all jobs of HI-criticality tasks** complete by their deadlines

Vestal. Preemptive scheduling of multi-criticality systems with varying degrees of execution time assurance. RTSS 2007

# Priority assignment schemes

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Baruah, Burns, and Davis. Response-time analysis for mixed-criticality systems. RTSS 2011

Vestal. Preemptive scheduling of multi-criticality systems with varying degrees of execution time assurance. RTSS 2007



**PERIOD TRANSFORMATION**

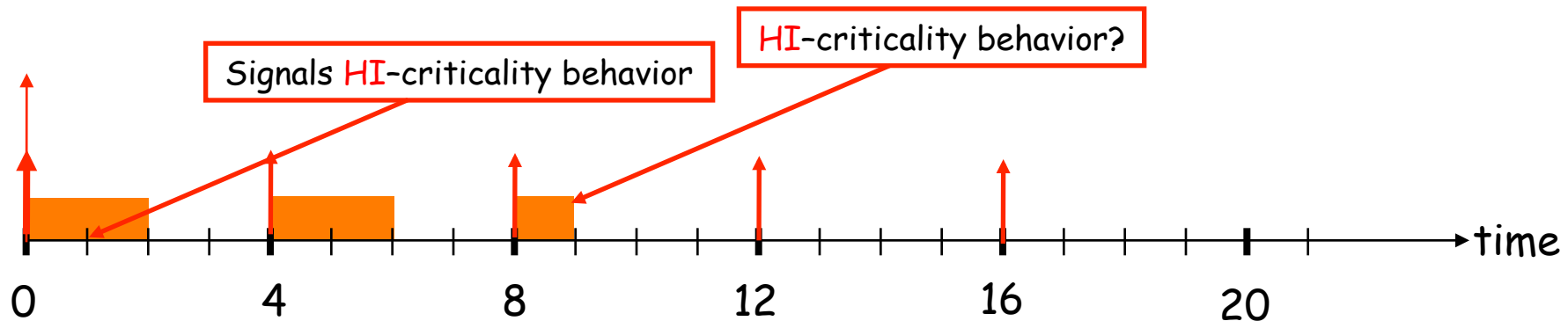
**VS**

**ADAPTIVE MIXED CRITICALITY**

**Period transformation (PT):** Scale down the parameters of HI-criticality tasks to make a DM priority assignment criticality-monotonic

	criticality	$C_i(\text{LO})$	$C_i(\text{HI})$	$D_i = T_i$
$\tau_1$	HI	<del>5</del> 1	<del>10</del> 2	<del>20</del> 4
$\tau_2$	LO	2	2	4

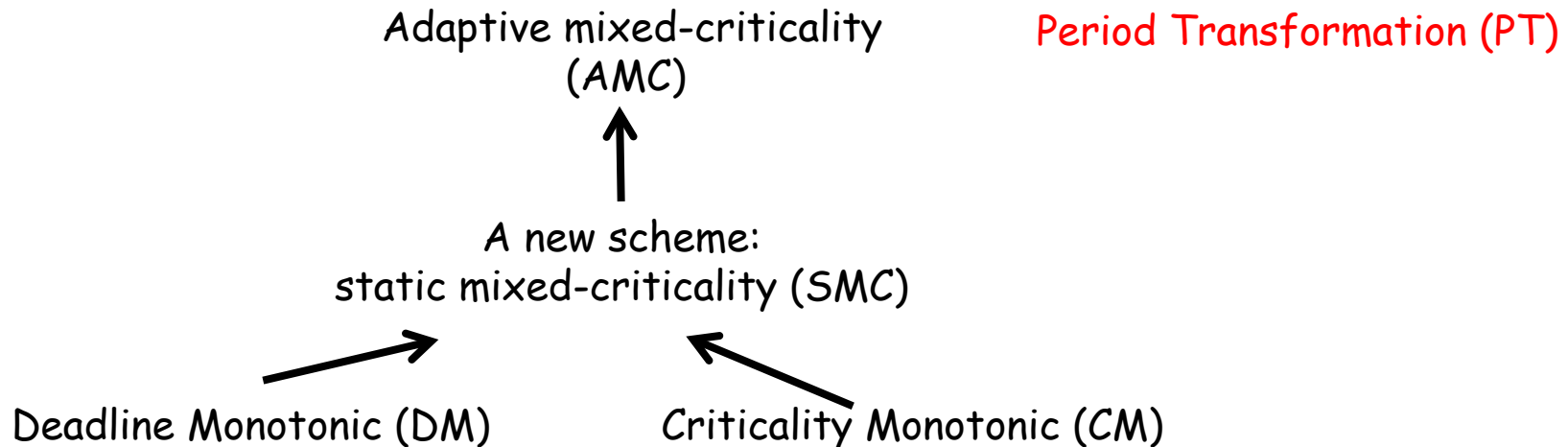
Scheduling the original task is not equivalent to scheduling a **mixed-criticality** task with the scaled-down parameters!!



# Observations - 1

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Result: AMC and PT are **incomparable**: each priority assignment scheme can schedule instances that the other cannot



# Observations - 2

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Result: AMC and PT are **incomparable**: each priority assignment scheme can schedule instances that the other cannot

PT on mixed-criticality instances

- Increase priority of **HI** tasks

- Improve **utilization bound** (as on non-MC tasks)

By making task periods harmonic

Result: AMC **dominates** PT for harmonic task systems

# Context and Conclusions

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**Correctness** and **efficiency** concerns in real-time systems

**Design methodologies** facilitate correctness

**Scheduling theory**: improve design methodologies to achieve efficiency

**The Fixed-Priority (RMA) design methodology**

- **platforms** increasingly **unpredictable** ➡ **mixed-criticality** implementations
- A **survey** of fixed-priority approaches to MC scheduling
- A **comparative evaluation** of the Period Transformation (PT) approach